Enhancing productivity on heterogeneous supercomputers with task-based programming model
Adrien Roussel, Mickael Boichot, Romain Pereira, Manuel Ferat

To cite this version:
Adrien Roussel, Mickael Boichot, Romain Pereira, Manuel Ferat. Enhancing productivity on heterogeneous supercomputers with task-based programming model. SIAM CSE 2023 - SIAM Conference on Computational Science and Engineering, Feb 2023, Amsterdam, Netherlands. cea-03994014

HAL Id: cea-03994014
https://cea.hal.science/cea-03994014
Submitted on 28 Feb 2023

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L’archive ouverte pluridisciplinaire HAL, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d’enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.
Enhancing Productivity on Heterogeneous Supercomputers with Task-based Programming Model

Adrien Roussel, Mickael Boichot, Manuel Ferat, Romain Pereira
CEA, DAM, DIF

SIAM CSE, Amsterdam
February 28, 2023
Introduction

- Most of the time, scientific applications are already parallelized for CPU (or partly for GPU).

- To take advantage of GPUs, developers have to rewrite their codes (in part? in whole?)

- Goals:
  - Build a preliminary analysis to determine the potential of the application for a GPU port.
  - Port mini-application to heterogeneous architectures with this information.

```c
int main(…)
{
    ...
}
```

Tedious & Time-consuming redesign
- data structures
- execution scheme
Contributions

Goal: Inform the user

- that there exists a test case of the application suitable for a GPU port
- about the strengths and weaknesses of the application for a GPU port

- We analyzed the variations of two metrics:
  - the degree of parallelism (→ saturate GPU cores)
  - and the working set (→ bounded GPU memory)

- Porting app. To GPU Computing
  - Using Task based prog. Model
    - MPI+OpenMP (tasks + target)

int main(...)
{
  /* parallel CPU code */
}

Parallel computations

working set

projection

Relate the two metrics

Extrapolate application behavior

GPU Porting

- varying dataset input size
Contents

1. Introduction

2. Application Characterization Method

3. Proxy Applications Analysis

4. Porting Proxy App to GPU: LULESH

5. Conclusion
2. Application Characterization Method
Targeted GPU Characterization

Maximizing FFMA throughput ("floating point fused multiple add")

\[ RP = \#\text{cores} \times L \]

where

- $L$ : 32 bits FFMA instruction latency

Characterization

**Targeted GPU**

- **L2 Cache**
- **DRAM**

**Required Parallelism (RP)**

**RAM (GB)**

Already known (GPU documentation)
Metrics to Characterize the Application

Applications analyzed at Loop Granularity (already parallelized with OpenMP)

L: for (i = 0; i < N; ++i)
    compute;

- Set of data necessary to run the loop
- Working Set measured at two Granularities, i.e., the size of smallest accessed memory element (using DynamoRIO)
- Parallelism Degree P(L)
- # independent Iterations (i.e., tripcount)
- $\text{WS}_{\text{malloc}}(L)$
- $\text{WS}_{\text{cache line}}(L)$
Extrapolation Methodology

Extract most \textit{time-consuming} loops

\begin{align*}
\forall L \in S, n \in \{n_1, n_2, \ldots\} & \rightarrow \text{measure} \\
\text{• working set } WS(L) & \\
\text{• parallelism degree } P(L)
\end{align*}

\begin{align*}
\forall L \in S, \quad WS(L) & = f(P(L)) \\
T & = g(n)
\end{align*}

<table>
<thead>
<tr>
<th>loop</th>
<th>\textit{working set}</th>
<th>\textit{parallelism degree}</th>
<th>input size</th>
</tr>
</thead>
<tbody>
<tr>
<td>L</td>
<td>WS(L)</td>
<td>P(L)</td>
<td>n_1</td>
</tr>
<tr>
<td>L</td>
<td>WS(L)</td>
<td>P(L)</td>
<td>n_2</td>
</tr>
<tr>
<td>...</td>
<td>...</td>
<td>...</td>
<td>...</td>
</tr>
</tbody>
</table>

For each loop

\textit{working set}

\textit{parallelism degree}

\textit{regression}
Extrapolation Methodology

Pre-selection:
- Extract most time-consuming loops

Measure:
- For all $L \in S$, $n \in \{n_1, n_2, \ldots\}$
- Working set $WS(L)$
- Parallelism degree $P(L)$

Inference:
- For all $L \in S$
  - $WS(L) = f(P(L))$
  - $P(L) = g(n)$

Projection:
- Test cases fitting target GPU memory and exposing enough parallelism
3. Proxy Applications Analysis

LULESH, miniFE
Experimental Protocol

► Evaluation Machine

<table>
<thead>
<tr>
<th></th>
<th>Intel Xeon E5-2680</th>
</tr>
</thead>
<tbody>
<tr>
<td>processor</td>
<td></td>
</tr>
<tr>
<td>#sockets</td>
<td>2</td>
</tr>
<tr>
<td>#cores/socket</td>
<td>8</td>
</tr>
<tr>
<td>RAM (GB)</td>
<td>128</td>
</tr>
<tr>
<td>OS</td>
<td>Linux 3.10</td>
</tr>
<tr>
<td>Compiler</td>
<td>Intel 20.0.0</td>
</tr>
</tbody>
</table>

► Evaluated mini-applications

- **LULESH** (hydrodynamic simulation)
- **MiniFE** (finite element problem)
- **Quicksilver** (Monte Carlo particle transport)
- **XSBench** (Monte Carlo particle transport kernel)

*All these applications have CPU OpenMP, GPU CUDA and HIP versions.*

► Considered targeted GPUs

<table>
<thead>
<tr>
<th>GPU</th>
<th>RAM</th>
<th>#cores</th>
<th>minimal required parallelism RP</th>
</tr>
</thead>
<tbody>
<tr>
<td>V100</td>
<td>16 GB</td>
<td>5120</td>
<td>20480</td>
</tr>
<tr>
<td>A100 40GB</td>
<td>40 GB</td>
<td>6912</td>
<td>27648</td>
</tr>
<tr>
<td>A100 80GB</td>
<td>80 GB</td>
<td>6912</td>
<td>27648</td>
</tr>
<tr>
<td>MI100</td>
<td>32 GB</td>
<td>7680</td>
<td>30720</td>
</tr>
</tbody>
</table>
Working set of LULESH selected loops as a function of their tripcount

Projection of dataset input size $s_{\text{max}}$ saturating GPU memory and corresponding $\rho$ ratio ($\rho(L) = P(L) / \text{RP}$)

- Sufficient exposed parallelism ($\rho \gg 1$)
- Limiting factor: memory size
- Advised data transfers granularity: malloc

- Similar conclusions for miniFE
4. Porting Proxy App to GPU: LULESH
Porting LULESH to MPI+OpenMP(tasks)

- LULESH
  - Hydrodynamic problems solver - Iterative simulation
  - Original code express loop-based parallelism (parallel-for)

- Porting to MPI+OpenMP (tasks)
  - Loops splitted into tasks with dependencies
  - Single producer/multi-consumer scheme
  - MPI communications nested within OpenMP tasks
    - Overlap communications by computations

```c
for (int j = 0 ; j < n_nodes; j += BS)
{
  # pragma omp task depend(out : mesh->nodes[j])
  {
    for (int i = j ; i < j + BS ; ++i)
      work_on_nodes(mesh->nodes[i]);
  }
}

for (int j = 0 ; j < n_cells; j += BS)
{
  # pragma omp task depend(in : [...])
  {
    for (int i = j ; i < j + BS ; ++i)
      work_on_cells(mesh->cells[i]);
  }
}
```

Feb. 28, 2023
Adrien Roussel - SIAM CSE
Concept Generalization for Heterogeneous Architectures with OpenMP Target

- **Porting to MPI+OpenMP (tasks + GPU target)**
  - GPU « super-tasks » merging several CPU tasks
  - From CPU tasks to GPU offloading while enhancing asynchronism
    - Sending/Retrieving data to/from GPU with ‘update’ clause
    - Executing GPU kernels through ‘nowait’ clause
  - Offloading **6 loops given by application characterization**
    - IntegrateStressForElems, CalcKinematicsForElems, CalcAccelerationForNodes, CalcFBHourglassForceForElems1, CalcHourglassControlForElems, CalcMonotonicQGradientsForElems
- **Use of cudaMallocHost for the GPU buffers**
  - Pinned memory to enable asynchronous execution

```c
1 for (int j = 0 ; j < n_nodes; j += BS) 
2 { 
3     # pragma omp task depend(out : mesh->nodes[j]) 
4     { 
5         for (int i = j ; i < j + BS ; ++i) 
6             work_on_nodes(mesh->nodes[i], buffer[i]); 
7     } 
8 } 
9 buffer = cudaMallocHost(…); 
[...]
```

```c
1 for (int j = 0 ; j < n_nodes; j += BS) 
2 { 
3     # pragma omp target update to(buffer[j:BS]) \ 
4       depend(out : mesh->nodes[j]) 
5     # pragma omp target 
6       teams distribute parallel for nowait \ 
7       depend(iterator(i=0:N_BS), out: mesh->nodes[j+i*BS]) 
8     { 
9         for (int i = j ; i < j + N_BS * BS ; ++i) 
10             work_on_nodes(mesh->nodes[i], buffer[i]); 
11     } 
12     # pragma omp target update from(buffer[j:BS]) \ 
13     depend(out: mesh->nodes[j]) 
14 } 
```

Adrien Roussel - SIAM CSE
# Experimental Protocol

<table>
<thead>
<tr>
<th>Processor</th>
<th>AMD EPYC 7H12</th>
</tr>
</thead>
<tbody>
<tr>
<td>#sockets</td>
<td>2</td>
</tr>
<tr>
<td>#cores/socket</td>
<td>64</td>
</tr>
<tr>
<td>GPU</td>
<td>NVIDIA A100</td>
</tr>
<tr>
<td>#GPU/socket</td>
<td>2</td>
</tr>
<tr>
<td>RAM (GB)</td>
<td>256</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>MPC-OpenMP</th>
<th>4.1.0</th>
</tr>
</thead>
<tbody>
<tr>
<td>(No Offloaded part)</td>
<td></td>
</tr>
<tr>
<td>LLVM</td>
<td>16.x</td>
</tr>
<tr>
<td>(Offloaded part)</td>
<td></td>
</tr>
<tr>
<td>Open MPI</td>
<td>4.0.5</td>
</tr>
</tbody>
</table>
LULESH Performance Evaluation

- **Problem size**
  - 264³ elements
  - 80% of the NUMA domain memory capacity
  - 20% of the GPU memory capacity
  - 128 iterations

- **Placement**
  - Per MPI Process:
    - 16 threads (on 16 cores)
    - + 1 GPU

- **Weak scaling**
  - Median runtime of 10 runs
  - Each version scale very well
  - **OpenMP tasks** version significantly improves performances due to **cache reuse**
  - Significant performance gain with **GPU offloading**

<table>
<thead>
<tr>
<th>Number of MPI ranks</th>
<th>Median Time (s)</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>parallel-for</td>
</tr>
<tr>
<td>1</td>
<td>565.84</td>
</tr>
<tr>
<td>8</td>
<td>572.45</td>
</tr>
<tr>
<td>27</td>
<td>581.26</td>
</tr>
</tbody>
</table>
5. Conclusion
Conclusion

- An approach: :

  \[
  \text{int main(...)}\{
  \ldots
  \}
  \mapsto \{ \text{test cases such that} \\
  \quad \text{int main(...)}\{
  \ldots
  \}
  \quad \text{fit on} \ \text{RAM} \\
  \quad \text{and do not lack parallelism} \\
  \}
  \]

- **complementary** to existing work (analysis of several application runs → extrapolation)

- We have **evaluated** our approach and shown that we can **qualitatively guide developer** in their GPU porting

- We have **ported** LULESH on heterogeneous machine with a **task-based programming model**
  - Significant gains are observed
Thank you!

Questions?